CE 530 Molecular Simulation

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Lecture 25 Efficiencies and Parallel Methods

David A. Kofke Department of Chemical Engineering SUNY Buffalo kofke@eng.buffalo.edu

Look-Up Tables

O Evaluation of interatomic potential can be time-consuming

• For example, consider the exp-6 potential

$$u(r) = -\frac{A}{r^6} + Be^{-Cr}$$

• Requires a square root and an exponential

O Simple idea:

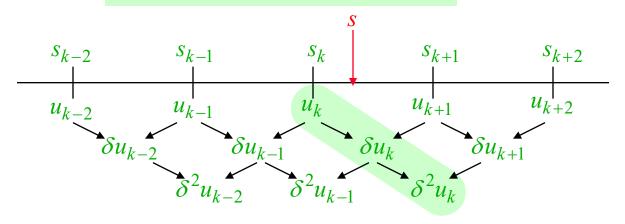
• Precompute a table of values at the beginning of the simulation and use it to evaluate the potential via interpolation

Interpolation

O Many interpolation schemes could be usedO e.g., Newton-Gregory forward difference method

- equally spaced values δs of $s = r^2$
- given $u_1 = u(s_1)$, $u_2 = u(s_2)$, etc.
- define first difference and second difference $\delta u_k = u_{k+1} - u_k$ $\delta^2 u_k = \delta u_{k+1} - \delta u_k$

• to get
$$u(s)$$
 for $s_k < s < s_{k+1}$, interpolate
$$u(s) \approx u_k + \xi \delta u_k + \frac{1}{2}\xi(\xi - 1)\delta^2 u_k \qquad \xi = (s - 1)\delta^2 u_k$$

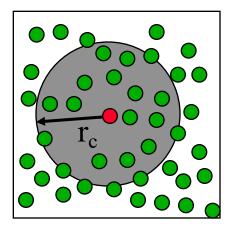


• forces, virial can be obtained using finite differences

 $(s_k)/\delta s$

Finding Neighbors Efficiently

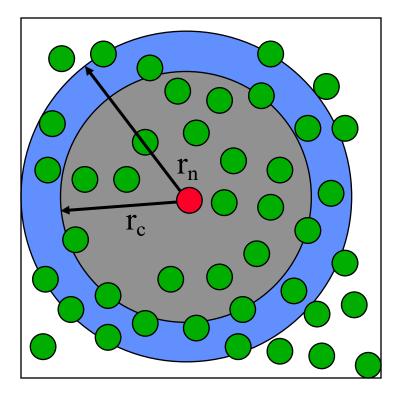
- O Evaluation of all pair interactions is an O(N²) calculation
- O Very expensive for large systemsO Not all interactions are relevant
 - potential attenuated or even truncated beyond some distance
- O Worthwhile to have efficient methods to locate neighbors of any molecule
- O Two common approaches
 - Verlet neighbor list
 - Cell list



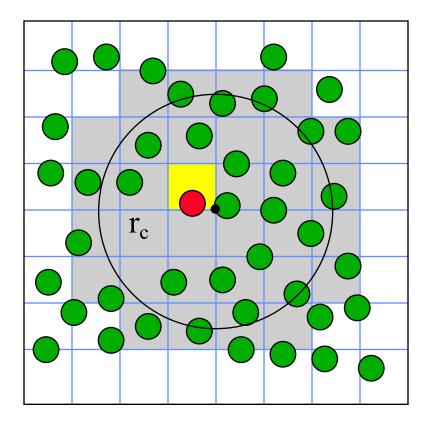
Verlet List

O Maintain a list of neighbors

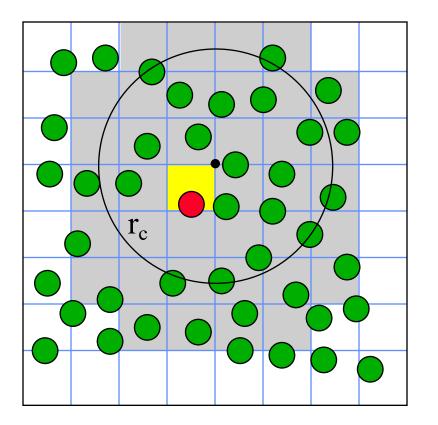
- Set neighbor cutoff radius as potential cutoff plus a "skin"
- O Update list whenever a molecule travels a distance greater than the skin thickness
- O Energy calculation is O(N)
- O Neighbor list update is $O(N^2)$
 - but done less frequently



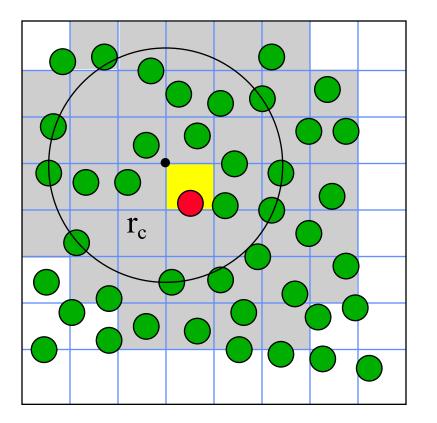
- O Partition volume into a set of cells
- O Each cell keeps a list of the atoms inside it
- At beginning of simulation set up neighbor list for each cell
 - list never needs updating



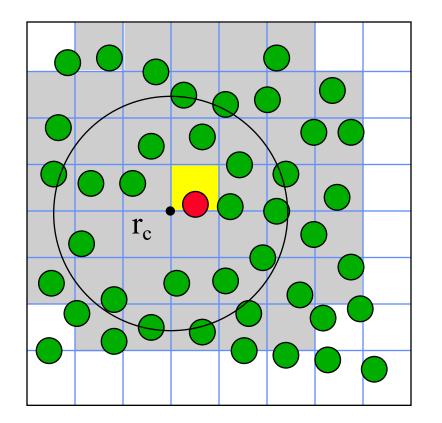
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- O Fewer unneeded pair interactions for smaller cells



Parallelizing Simulation Codes

O Two parallelization strategies

• Domain decomposition

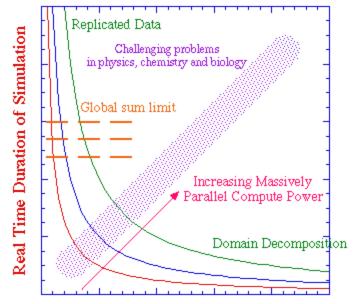
Each processor focuses on fixed region of simulation space (cell)Communication needed only with adjacent-cell processorsEnables simulation of very large systems for short times

• *Replicated data*

Each processor takes some part in advancing all molecules Communication among all processors required Enables simulation of small systems for longer times

Limitations on Parallel Algorithms

• Straightforward application of raw parallel power insufficient to probe most interesting phenomena



Number of Atomic Units

• Advances in theory and technique needed to enable simulation of large systems over long times

Figure from P.T. Cummings

Parallelizing Monte Carlo

- O Parallel moves in independent regions
 - moves and range of interactions cannot span large distances

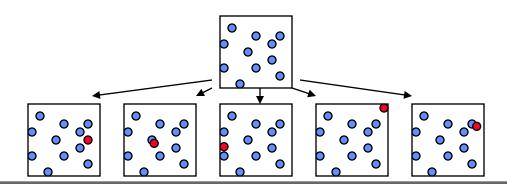
O Hybrid Monte Carlo

• *apply MC as bad MD, and apply MD parallel methods* time information lost while introducing limitations of MD

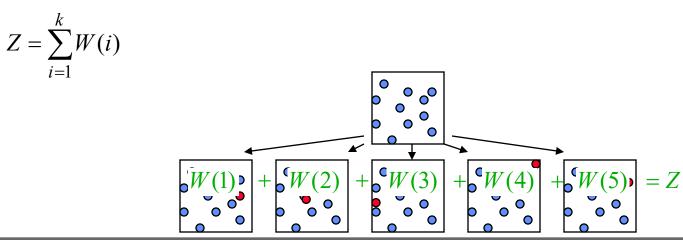
O Farming of independent tasks or simulations

- equilibration phase is sequential
- often a not-too-bad approach
- O Parallel trials with coupled acceptance
 - "Esselink" method

- O 1. Generate k trials from the present configuration
 - each trial handled by a different processor
 - useful if trials difficult to generate (e.g., chain configurational bias)

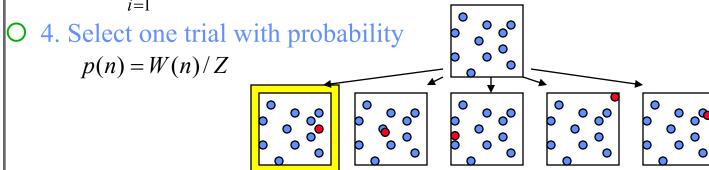


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 - e.g., Rosenbluth weight if CCB
 - more simply, W(i) = exp[-U(i)/kT]
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 - Need to evaluate a probability it would be generated from trial configuration
 - *Plan: Choose a path that includes the other (ignored) trials*



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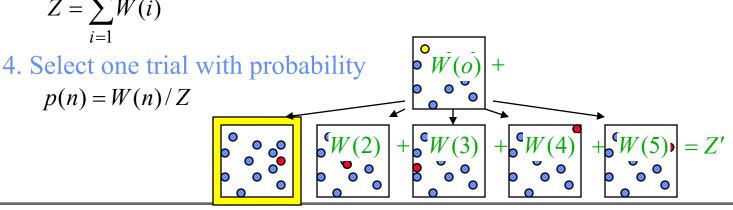
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 $Z' = Z - W(n) + W(o) \equiv R + W(o)$

O 8. Accept new trial with probability

$$p_{acc} = \min\left[1, \frac{Z}{Z'}\right]$$

Some Results

• Use of an incorrect acceptance probability

$$p_{acc} = \min\left[1, \frac{W(n)}{W(o)}\right]$$

$$p_{acc} = \min\left[1, \frac{W(n) + R}{W(o) + R}\right]$$

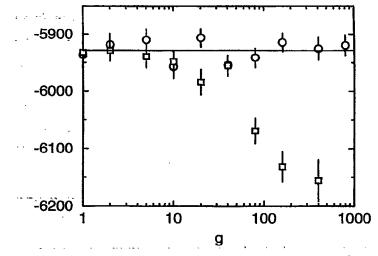
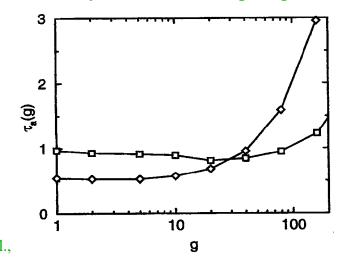


FIG. 1. Comparison of the total energy of pentane in silicalite as calculated from the sampling without the correction for the blas (\Box) with the correct sampling scheme (°). g is the number of chains grown in parallel and f = 1. The horizontal line is the average energy as calculated from the correct results.

O A sequential implemention

- O Average cpu time to acceptance of a trial
 - independent of number of trials g up to about g = 10
 - indicates parallel implementation with g = 10 would have 10 × speedup
 - *larger* g *is wasteful because acceptable configurations are rejected* only one can be accepted per move



Esselink et al., PRE, 51(2)

1995

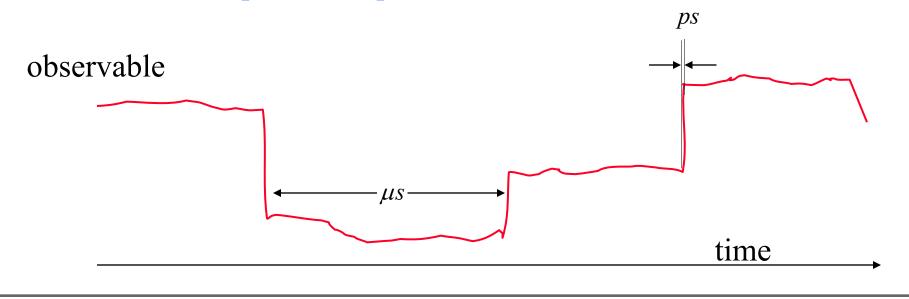
FIG. 2. Average time to acceptance $\tau_a(g)$ for methane (\diamond) and pentane (\Box) in silicalite for varying number of molecules g placed in parallel. The data are taken from Tables I and II.

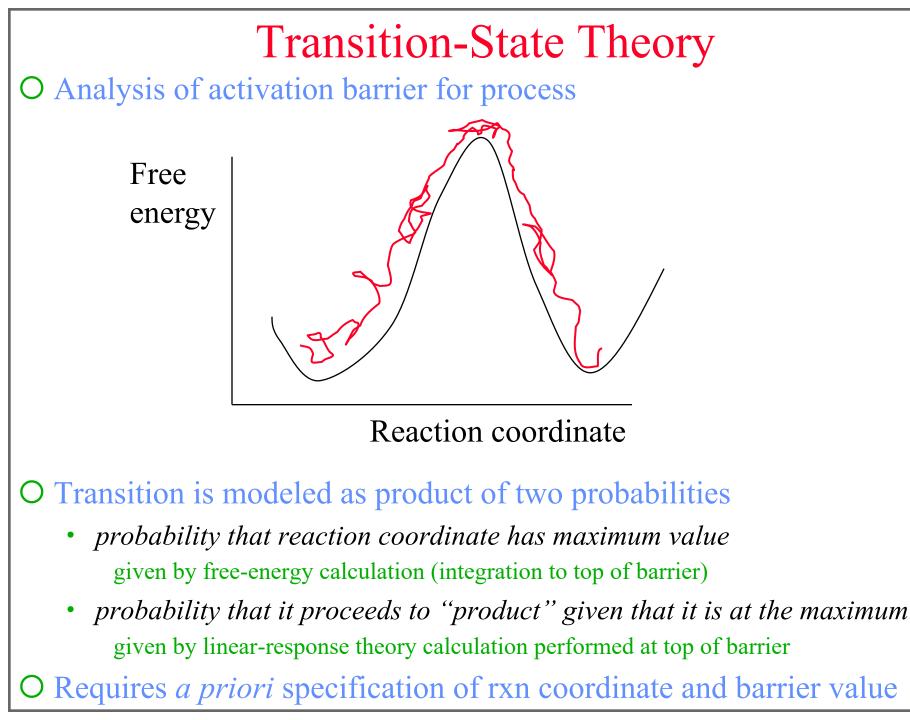
Simulation of Infrequent Events

O Some processes occur quickly but infrequently; e.g.

- rotational isomerization
- diffusion in a solid
- chemical reaction

O Time between events may be microseconds or longer, but event transpires over picoseconds





Parallel Replica Method (Voter's Method)

- O Establish several configurations with same coordinates, but different initial momenta
- O Specify criterion for departure from current "basin" in phase space
 - *e.g., location of energy minimum* evaluate with steepest-descent or conjugate-gradient methods
- O Perform simulation dynamics in parallel for different initial systems
- O Continue simulations until one of the replicas is observed to depart its local basin
- O Advance simulation clock by sum of simulation times of all replicas
- O Repeat beginning all replicas with coordinates of escaping replica

Theory Behind Voter's Method

• Assumes independent, uncorrelated crossing events • Probability distribution for a crossing event (sequential calculation)

 $p(t) = ke^{-kt}$ k = crossing rate constant

O Probability distribution for crossing event in any of M simulations

 $p_M(t) = \sum_{j=1}^{M} \left(\begin{array}{c} \text{probability of crossing} \\ \text{in simulation j at time t}_j \end{array} \right) \times \prod_{m=1}^{M} \left(\begin{array}{c} \text{probability simulation m} \\ \text{has yet had crossing event} \end{array} \right)$ $= \sum_{j=1}^{M} p(t_j) \times \prod_{m \neq j}^{M} \overline{p}(t_m)$ O No-crossing cumulative probability

$$\overline{p}(t) = \int_{t}^{\infty} p(\tau) d\tau = e^{-kt}$$

O Thus

$$p_M(t) = \sum_{j=1}^{M} k e^{-kt_j} \times \prod_{m \neq j}^{M} e^{-kt_m}$$
$$= M k e^{-kt_{sum}}$$

Rate constant for independent crossings is same as for individual crossings, if t_{sum} is used

Appealing Features of Voter's Method

O No *a priori* specification of transition path / reaction coordinate required

- Works well with multiple (unidentified) transition states
- *Does not require specification of "product" states*

O Parallelizes time calculation

• "Discarded" simulations are not wasted

O Works well in a heterogeneous environment of computing platforms

- OK to have processors with different computing power
- Parallelization is loosely coupled

Summary

O Some simple efficiencies to apply to simulations

- Table look-up of potentials
- Cell lists for identifying neighbors

O Parallelization methods

- Time parallelization is difficult
- Esselink method for parallelization of MC trials
- *Voter's method for parallelization of MD simulation of rare events*